Deductive Databases and Logic Programming (Winter 2003/2004)

Chapter 5: Practical Prolog Programming

- The Cut and Related Constructs
- Prolog vs. Pascal
- Definite Clause Grammars



After completing this chapter, you should be able to:

- explain the effect of the cut.
- write Prolog programs for practical applications.
- use context-free grammars in Prolog.

Overview

1. The Cut and Related Constructs

2. Prolog vs. Pascal

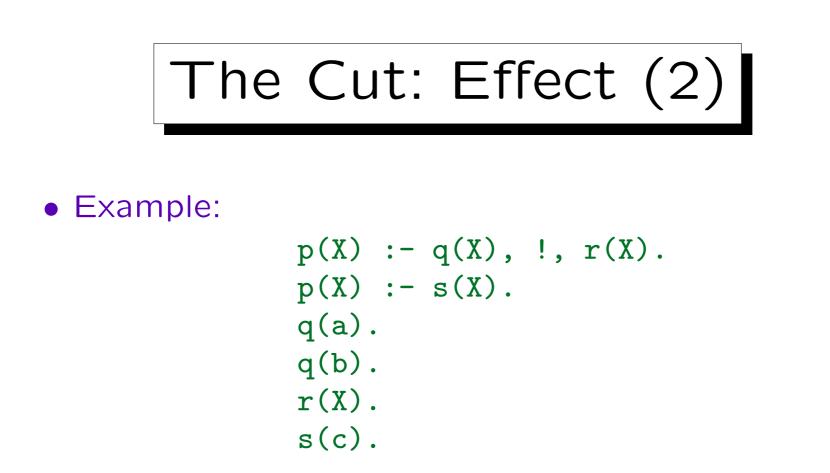
3. Definite Clause Grammars



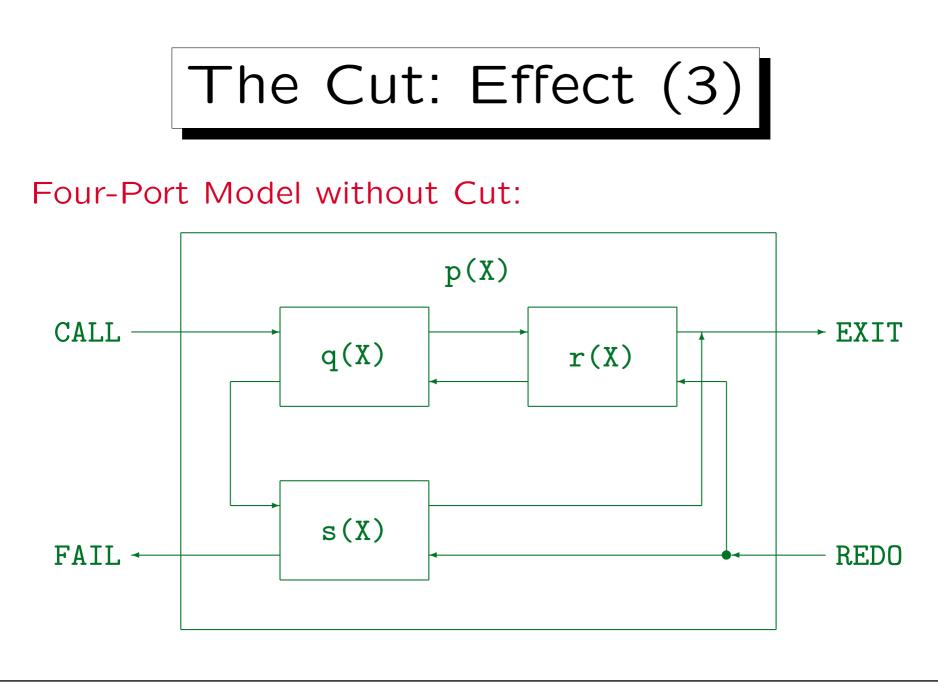
• The cut, written "!" in Prolog, removes alternatives that otherwise would have been tried during backtracking. E.g. consider this rule:

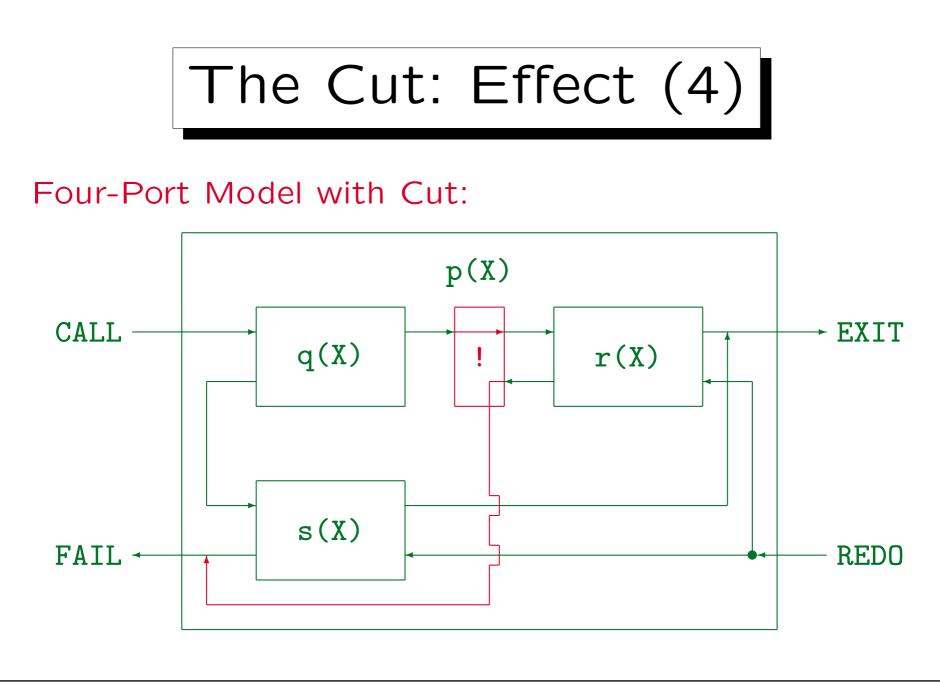
$$p(t_1,\ldots,t_k) := A_1, \ldots, A_m, !, B_1, \ldots, B_n.$$

- Until the cut is executed, processing is as usual.
- When the cut is reached, all previous alternatives for this call to the predicate p are removed:
 - $\diamond~$ No other rule about p will be tried.
 - \diamond No other solutions to the literals A, \ldots, A_m will be considered.



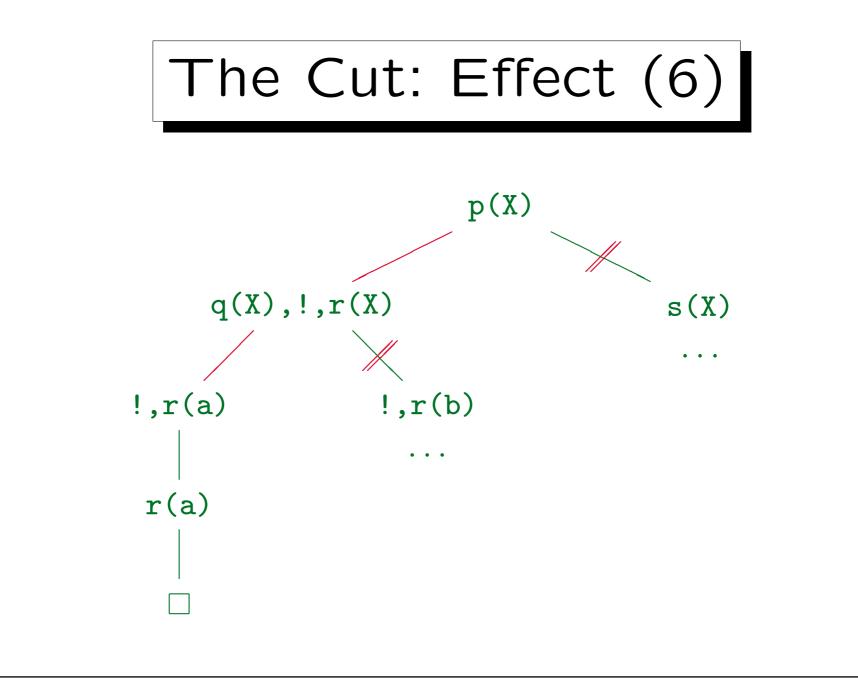
- With the cut, the query p(X) returns only X=a.
- Without the cut, the solutions are X=a, X=b, X=c.
- Exercise: Can the second rule about p ever be used?

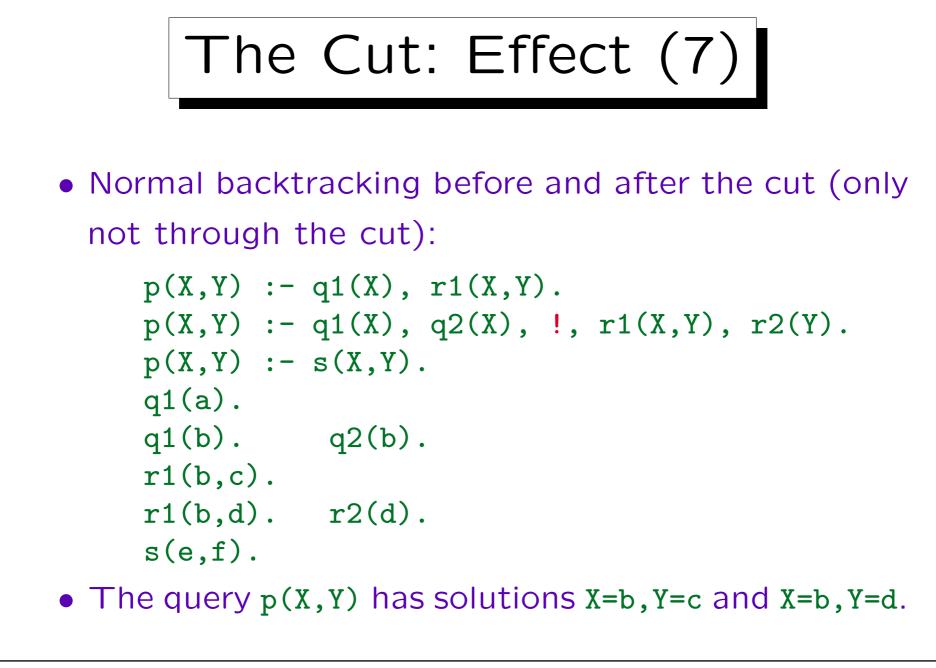


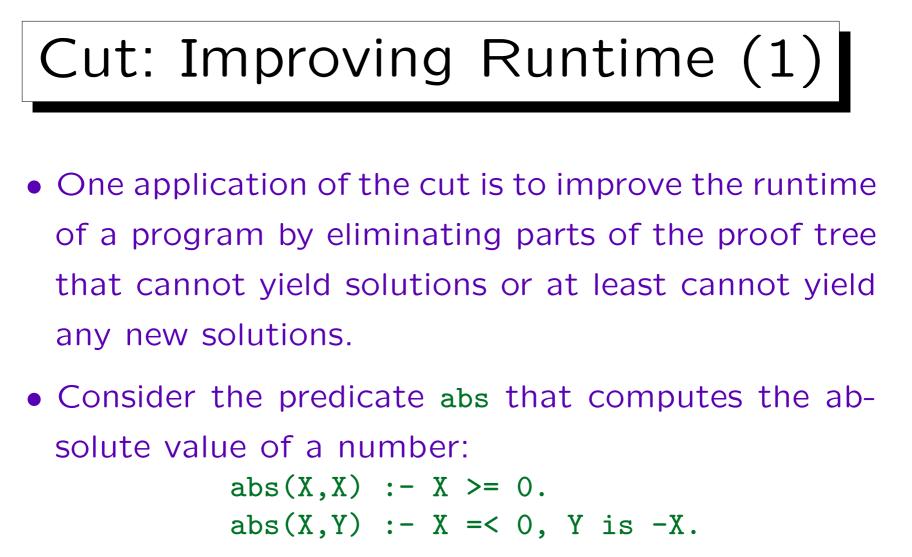




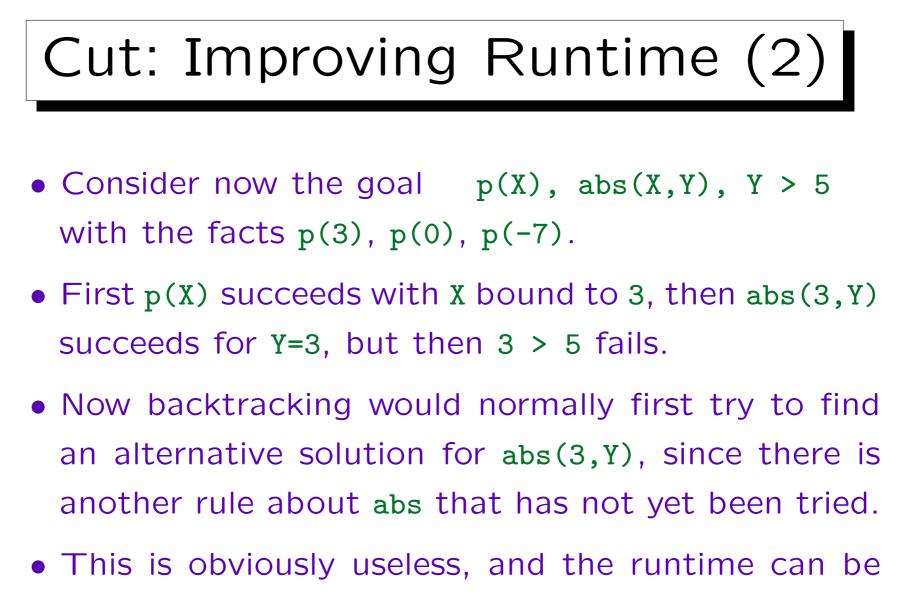
- A call to the cut immediately succeeds (like true).
- Any try to redo the cut not only fails, but immediately fails the entire predicate call.
- In the SLD-tree, the cut "cuts away" all still open branches between
 - the node where the cut was introduced (i.e. the child of which contains the cut), and
 - \diamond the node where the cut is the selected literal.



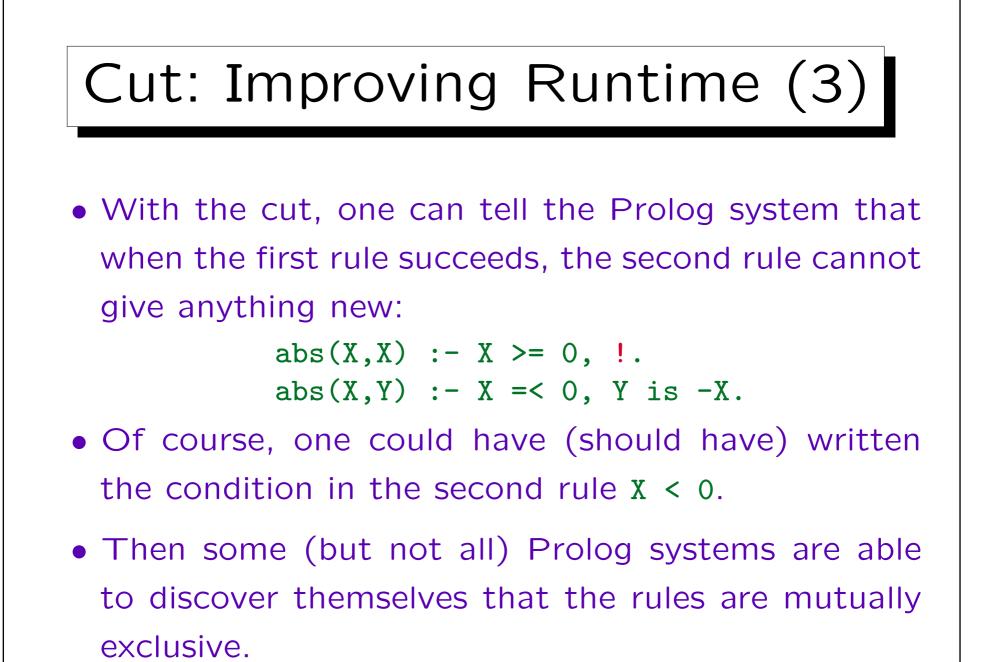




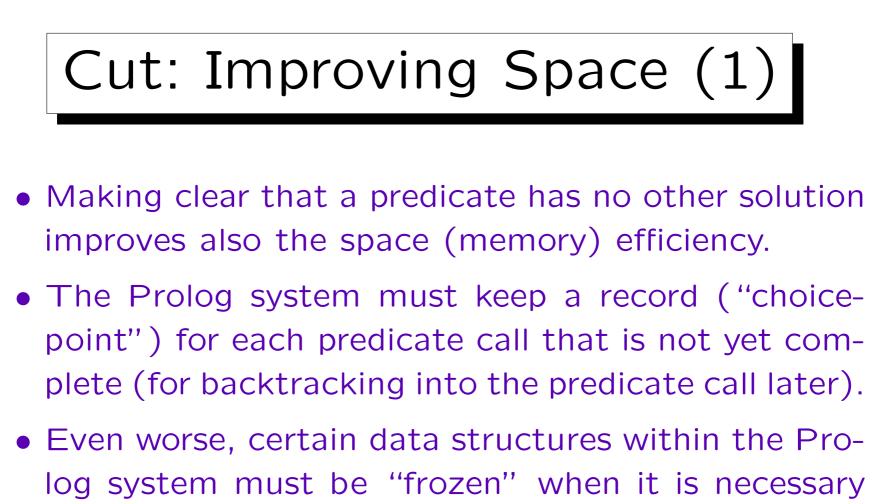
• When the first rule is successful, it is clear that the second rule does not have to be tried.



improved by immediately backtracking to p(X).



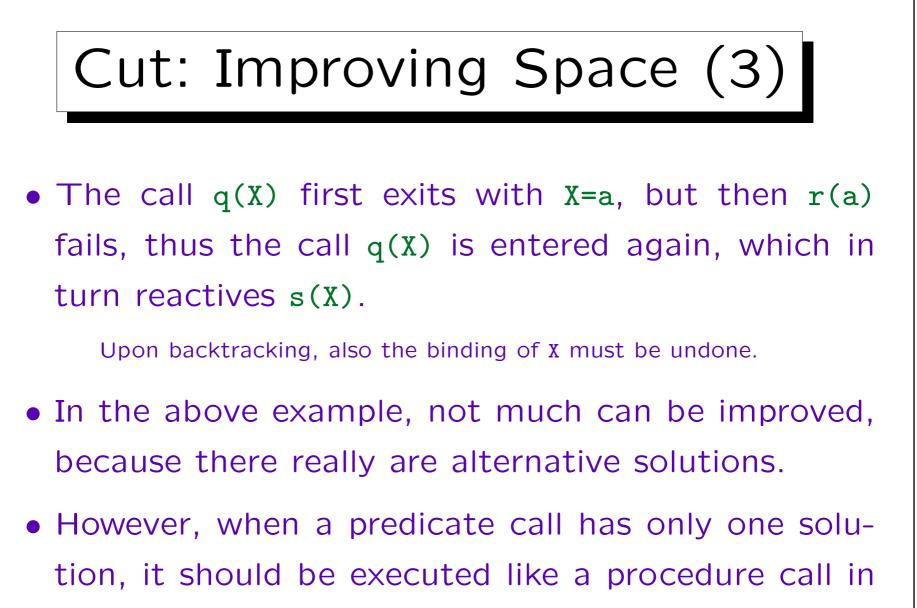
Stefan Brass: Deductive Databases and Logic Programming Univ. Halle, 2003



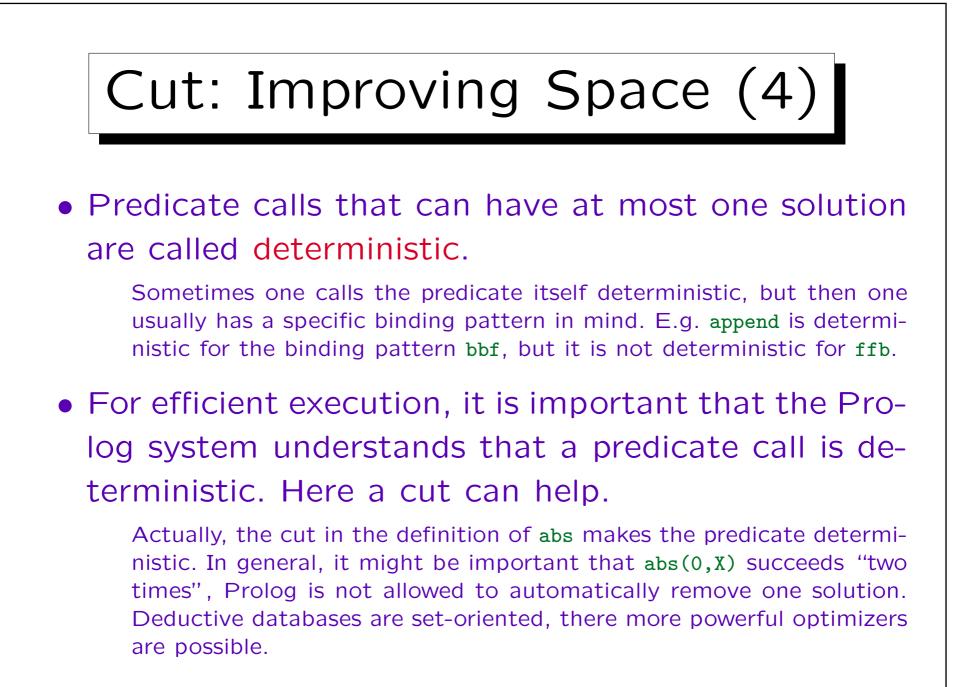
- to support later backtracking to this state.
- Then e.g. variable bindings must be logged (on the "trail") so that they can later be undone.

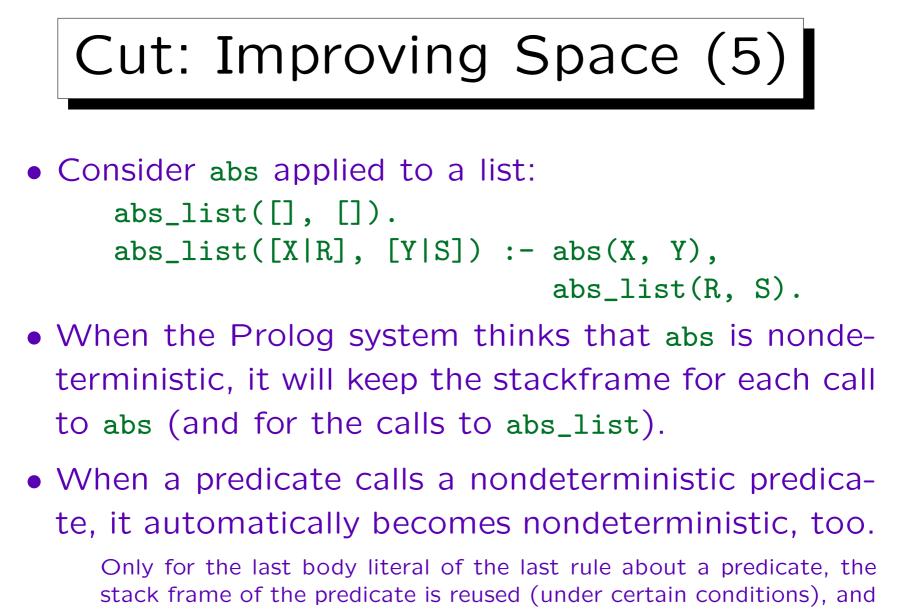


- In imperative languages, when a procedure call returns, its stack frame (containing local variables and other information) can be reused.
- In Prolog, this is not always the case, because it might be necessary to reactivate the procedure call and search for another solution.
- E.g. consider the following program:

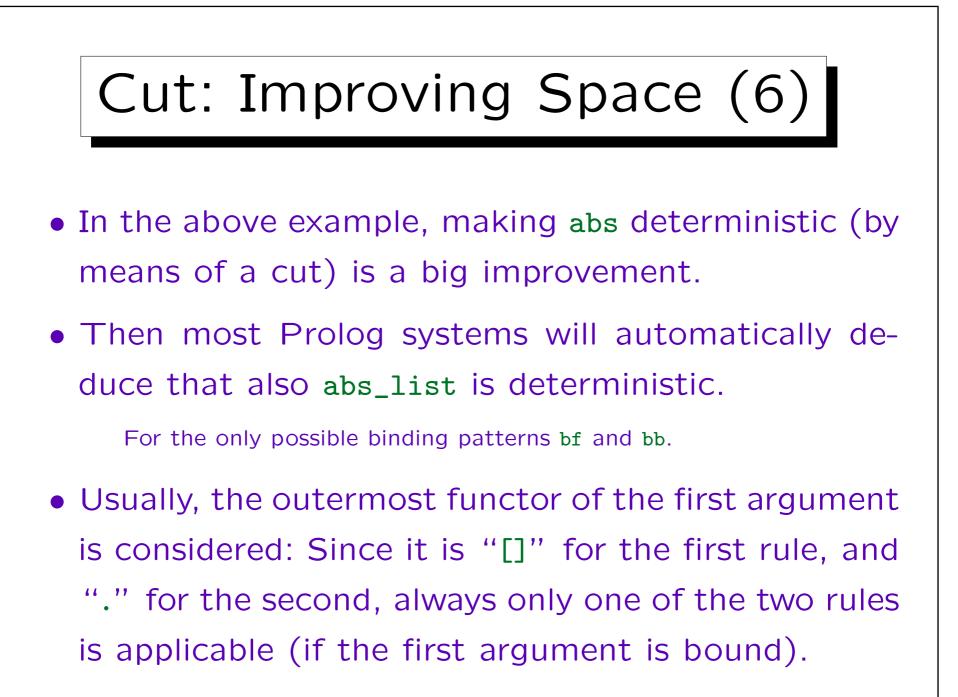


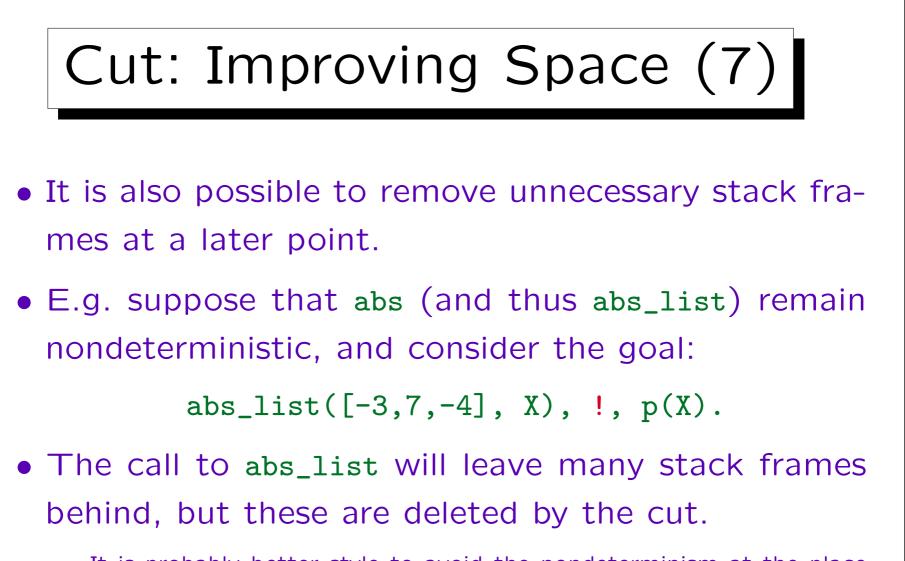
an imperative language.



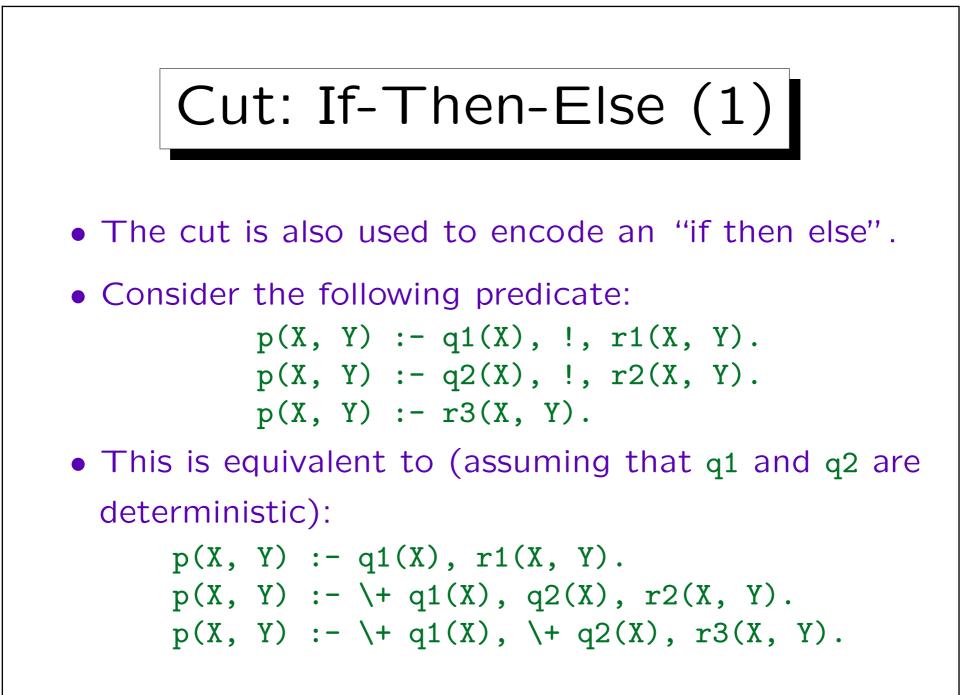


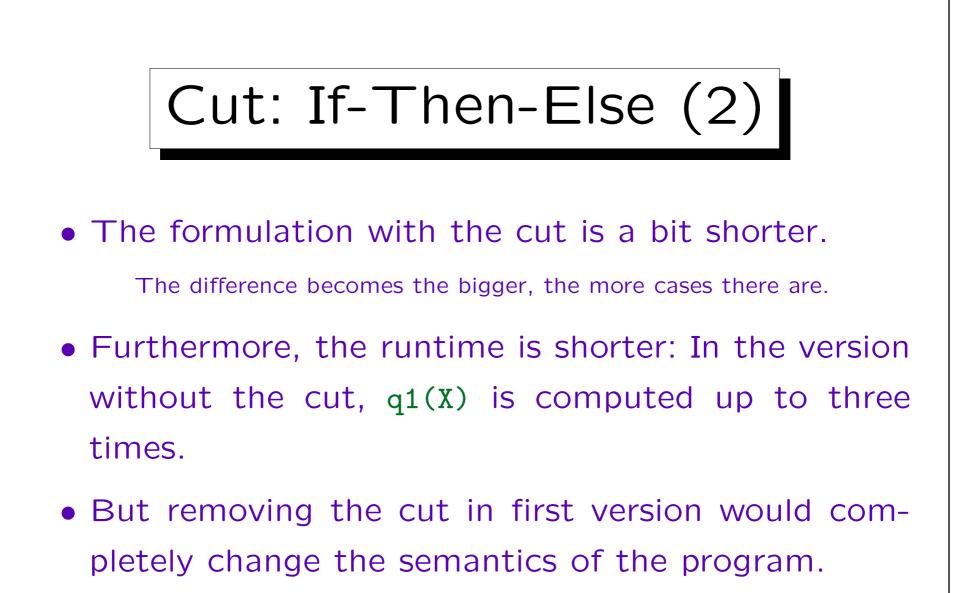
thus does not remain, even when this body literal is non-deterministic.



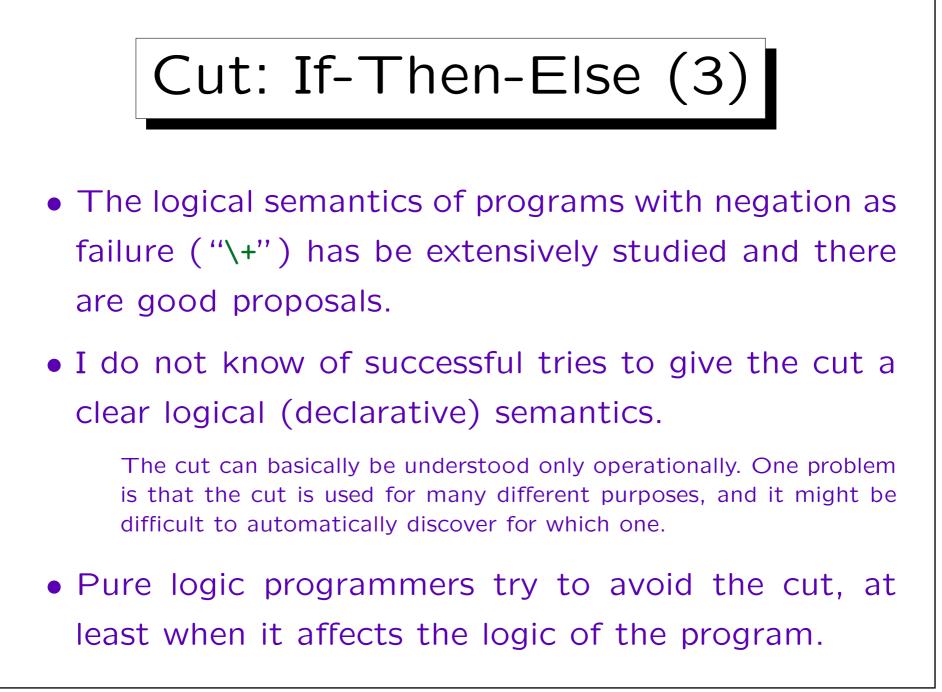


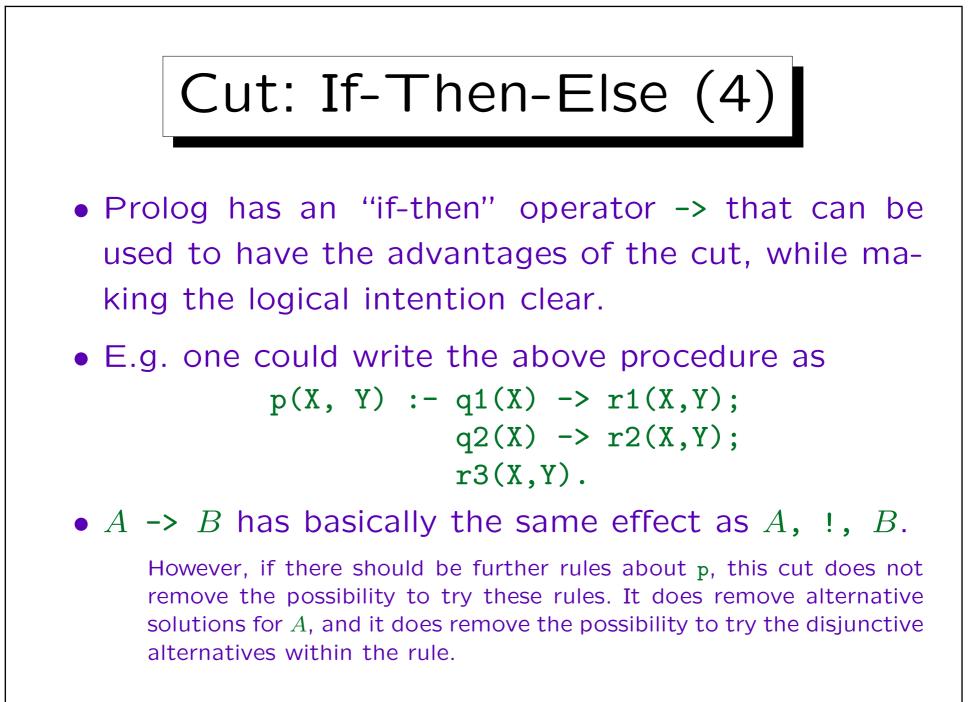
It is probably better style to avoid the nondeterminism at the place where it occurs. However, one should not use too many cuts, and it might be easier to clean up the stack only at a few places.

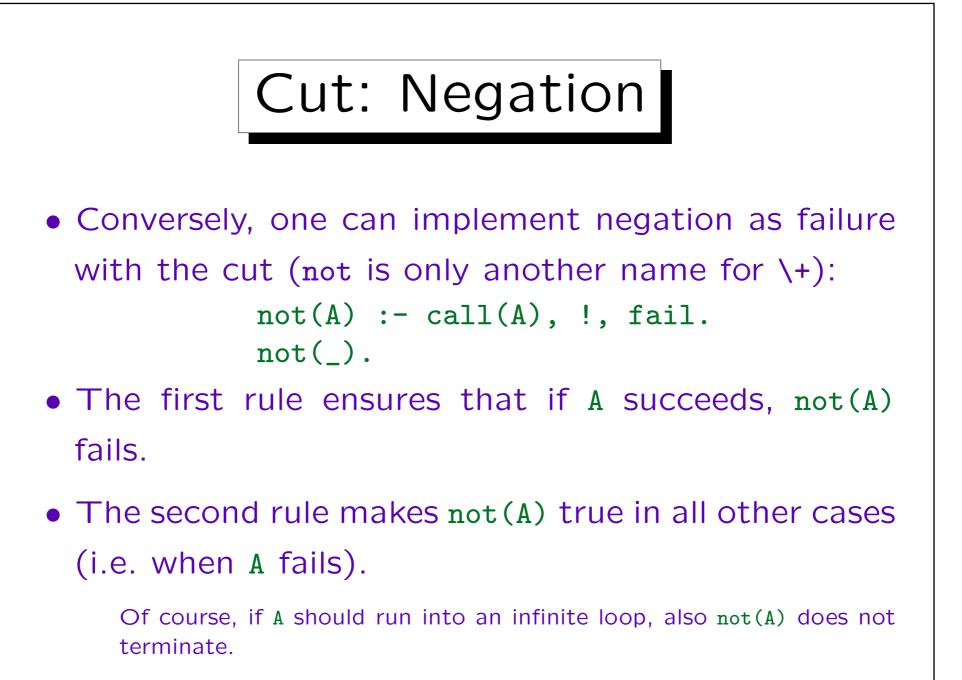


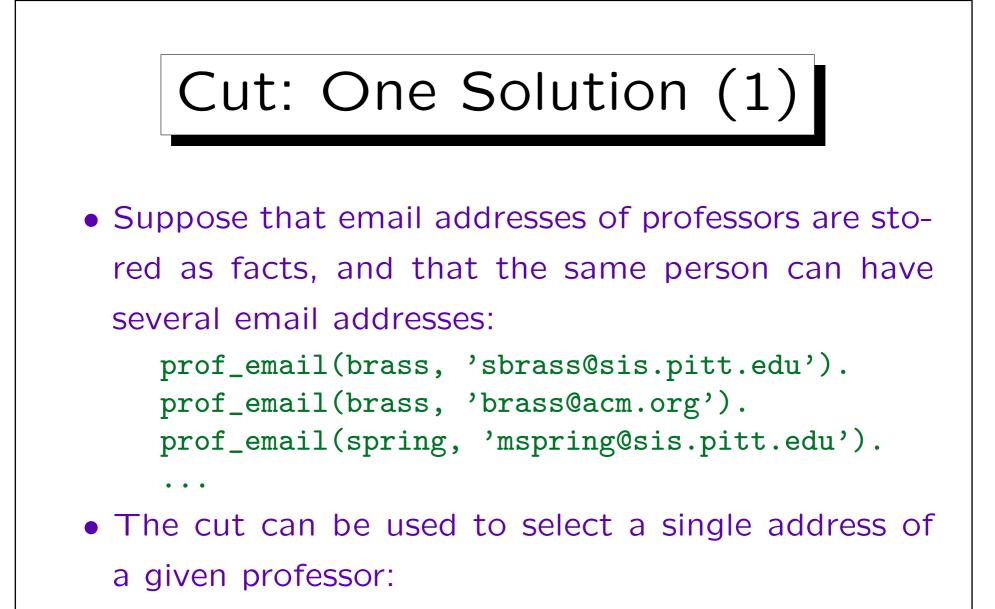


The cut is no longer only an "optimizer hint".

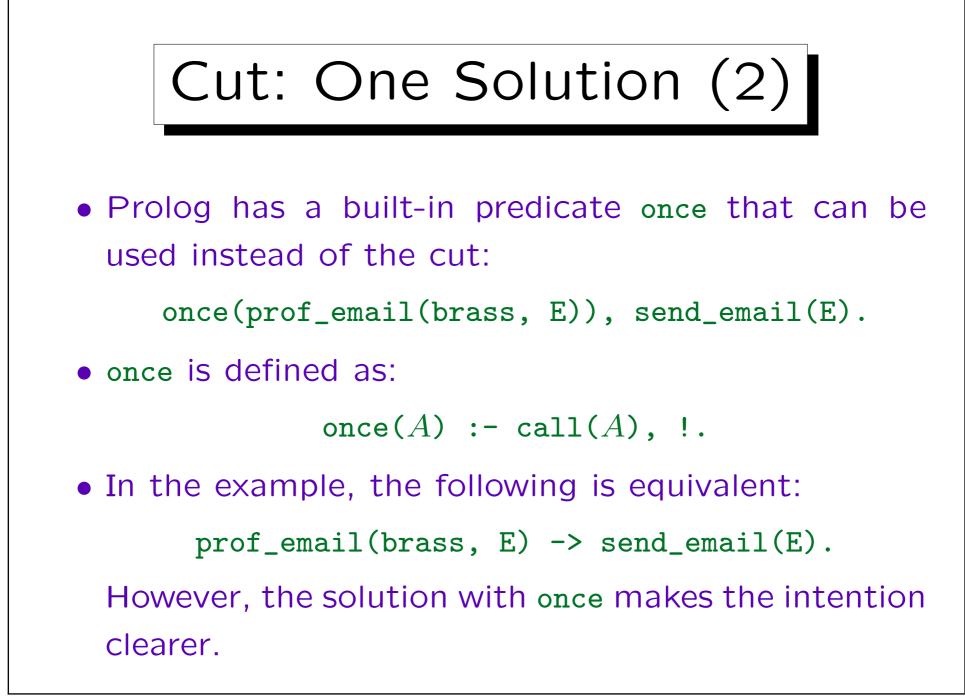


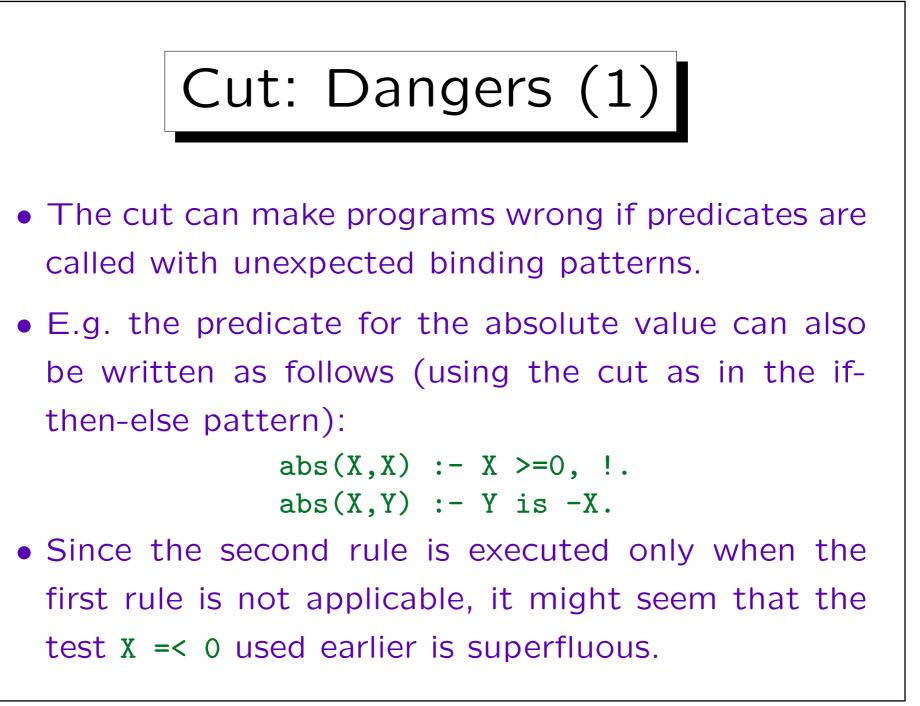






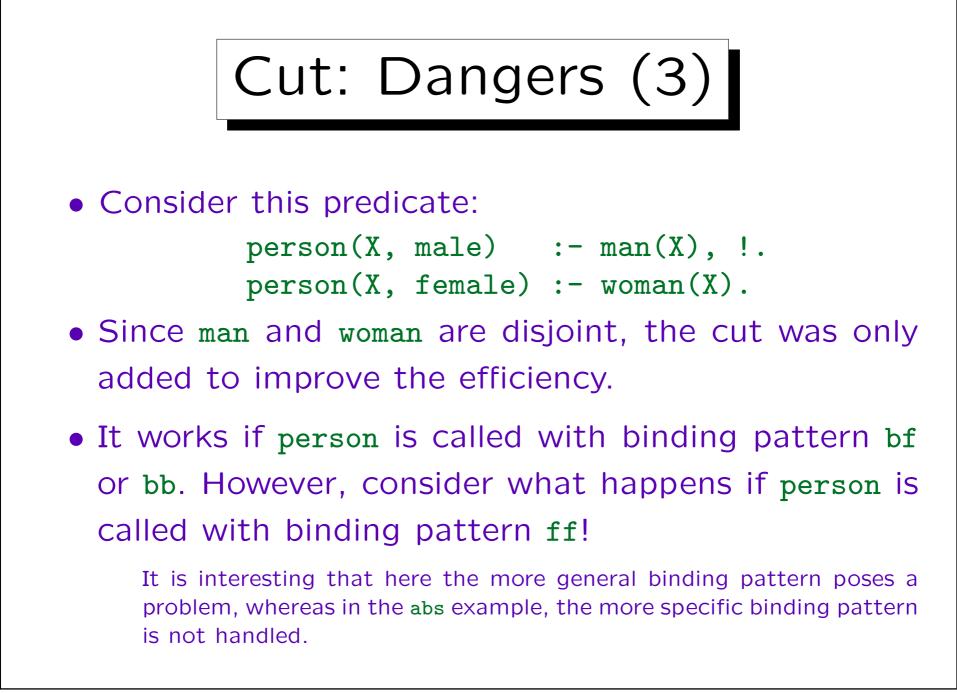
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prof_email(brass, E), !, send_email(E).
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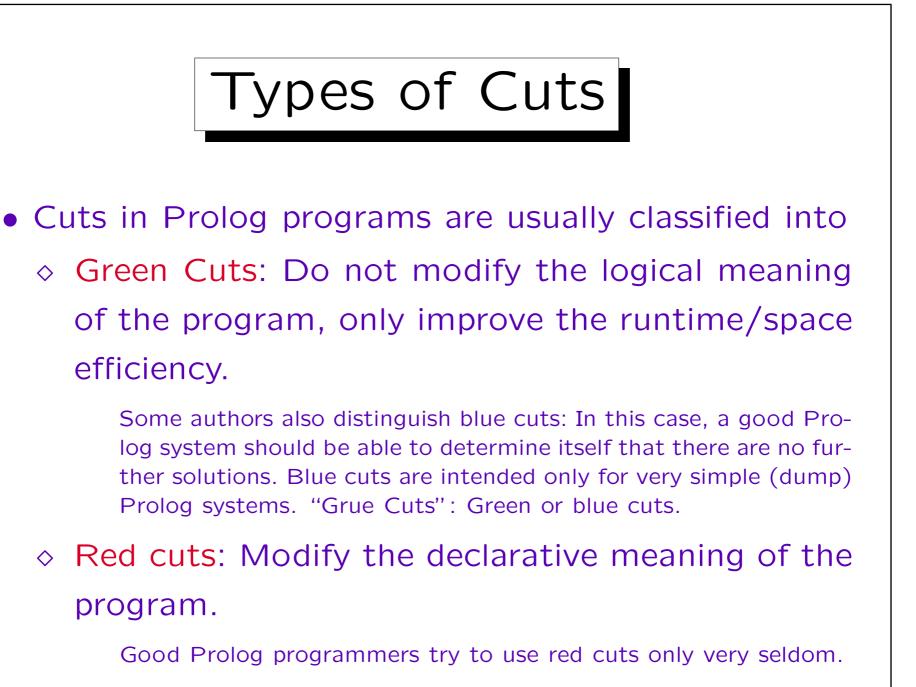


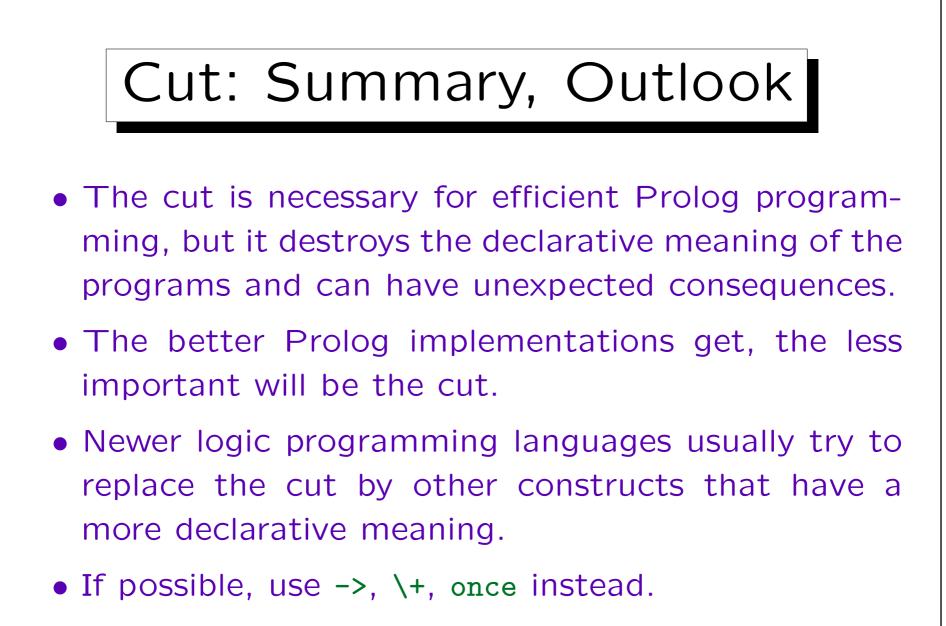


Cut: Dangers (2)

- This is indeed true for the binding pattern bf, but consider now the call abs(3,-3)!
- In general, the rule is that the cut must be exactly at the point where it is clear that this is the right rule: Not too early and not too late.
- Here the unification must happen after the cut: abs(X,Y) :- X >= 0, !, X = Y. abs(X,Y) :- Y is -X.
- This would work also with binding pattern bb.







• Use the cut only as last resort.